Lecture 25: Geodesic Paths

COMPSCI/MATH 290-04

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4/14/2016

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- Group Assignment 3 Out: First Deadline Monday 4/18.
 Final Deadline Tuesday 4/26
- ▷ Final Project Final Deadline 5/3 5:00 PM

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► Geodesics

- Dijkstra's / Fast Marching
- ▷ G2 Geodesic Histograms

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Geodesic Paths



Euclidean Path (shortest path of flying fly)

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Geodesic Paths



Euclidean Path (shortest path of flying fly)



Geodesic Path (shortest path of crawling ant)



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Geodesic paths on spheres lie along great circles



- Geodesic paths on spheres lie along great circles
- > Geodesic distance is the shortest geodesic path

SQC.



- Geodesic paths on spheres lie along great circles
- > Geodesic distance is the shortest geodesic path
- ▷ What is the geodesic distance between two points \vec{P} and \vec{Q} on a sphere centered at the origin with radius R?

SQC.

What is the geodesic distance between two points \vec{P} and \vec{Q} on a sphere centered at the origin with radius *R*?

$$R\cos^{-1}\left(\frac{\vec{P}\cdot\vec{Q}}{||\vec{P}||||\vec{Q}||}\right) = R\cos^{-1}\left(\frac{\vec{P}\cdot\vec{Q}}{R^2}\right)$$

Remember SLERP??

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Another Geodesic Mesh Example



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- \triangleright Geodesics
- Dijkstra's / Fast Marching
- ▷ G2 Geodesic Histograms

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Dijkstra's Algorithm Review

```
def Dijsktra(Graph, source):
 list dists
 list prev
 dist[source] = 0
 Queue O
 for vertex v in Graph:
     if v not source:
         dists[v] = Infinity
         prev[v] = Undefined
     O.add(v, dists[v])
 while len(Q) > 0:
     u = Q.qetMin()
     for v in neighbors(u):
         d = dists[u] + length(u, v)
         if d < dists[v]:
             dists[v] = d
             prev[v] = u
             Q.decreasePriority(v, d)
 return (dist, prev)
```

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Dijkstra's Algorithm Review

```
def Dijsktra(Graph, source):
 list dists
                                             What is the
 list prev
                                             worst case
 dist[source] = 0
                                             behavior for
 Queue O
 for vertex v in Graph:
                                               \triangleright V ver-
     if v not source:
                                                 tices
          dists[v] = Infinity
                                               \triangleright E edges
          prev[v] = Undefined
     O.add(v, dists[v])
                                             for a bal-
 while len(Q) > 0:
                                                      min
                                             anced
     u = Q.qetMin()
                                             heap Q?
     for v in neighbors(u):
          d = dists[u] + length(u, v)
          if d < dists[v]:
              dists[v] = d
              prev[v] = u
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 return (dist, prev)
```

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Dijkstra's Algorithm Review

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                                             heap Q?
     for v in neighbors(u):
          d = dists[u] + length(u, v)
          if d < dists[v]:
                                             O((E+V)\log(V))
              dists[v] = d
              prev[v] = u
              Q.decreasePriority(v, d)
 return (dist, prev)
```

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8x8 Cartesian Grid: Side Length 1



Shortest path along mesh is length $7\sqrt{2}$

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8x8 Cartesian Grid: Side Length 1



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8x8 Cartesian Grid: Side Length 1



Shortest path along mesh is 14

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Does refining the grid help? 15x15 Cartesian Grid: Side Length 0.5



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Does refining the grid help? 15x15 Cartesian Grid: Side Length 0.5



Nope!

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In general, mesh biases the solution!



Fast Marching

A modification of Dijkstra's algorithm to cut through triangles



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Fast Marching

A modification of Dijkstra's algorithm to cut through triangles



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- \triangleright Geodesics
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Mesh Isomorphisms

An isomorphism preserves all pairwise geodesic distances





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Mesh Isomorphisms

Contrast with Euclidean





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