# Lecture 25: Geodesic Paths 

## COMPSCI/MATH 290-04

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4/14/2016

## Announcements

$\triangleright$ Group Assignment 3 Out: First Deadline Monday 4/18. Final Deadline Tuesday 4/26
$\triangleright$ Final Project Final Deadline 5/3 5:00 PM

## Table of Contents

- Geodesics
$\triangleright$ Dijkstra's / Fast Marching
$\triangleright$ G2 Geodesic Histograms


## Geodesic Paths



## Euclidean Path (shortest path of flying fly)

## Geodesic Paths




Geodesic Path (shortest path of crawling ant)

## Geodesic Paths on Spheres




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$\triangleright$ Geodesic paths on spheres lie along great circles

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$\triangleright$ Geodesic distance is the shortest geodesic path
$\triangleright$ What is the geodesic distance between two points $\vec{P}$ and $\vec{Q}$ on a sphere centered at the origin with radius $R$ ?

## Geodesic Paths on Spheres

What is the geodesic distance between two points $\vec{P}$ and $\vec{Q}$ on a sphere centered at the origin with radius $R$ ?

$$
R \cos ^{-1}\left(\frac{\vec{P} \cdot \vec{Q}}{\|\vec{P}\|\|\vec{Q}\|}\right)=R \cos ^{-1}\left(\frac{\vec{P} \cdot \vec{Q}}{R^{2}}\right)
$$

Remember SLERP??

## Another Geodesic Mesh Example




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$\triangleright$ Geodesics

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## Dijkstra's Algorithm Review

```
def Dijsktra(Graph, source):
```

    list dists
    list prev
    dist[source] \(=0\)
    Queue Q
    for vertex \(v\) in Graph:
        if \(v\) not source:
            dists[v] = Infinity
            prev[v] = Undefined
        Q.add(v, dists[v])
    while len (Q) > 0 :
$u=Q . g e t M i n()$
for $v$ in neighbors(u):
$\mathrm{d}=$ dists[u] + length(u, $v)$
if $d<d i s t s[v]:$
dists[v] = d
prev[v] = u
Q.decreasePriority (v, d)
return (dist, prev)

## Dijkstra's Algorithm Review

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        Q.add(v, dists[v])
while len(Q) > 0:
    u = Q.getMin()
    for v in neighbors(u):
        d = dists[u] + length(u, v)
        if d < dists[v]:
            dists[v] = d
            prev[v] = u
            Q.decreasePriority(v, d)
return (dist, prev)
```


## What is the worst case behavior for

$\triangleright V$ vertices
$\triangleright E$ edges
for a balanced min heap $Q$ ?
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## What is the worst case behavior for

 tices
$\triangleright E$ edges
for a balanced min heap $Q$ ?

$$
O((E+V) \log (V))
$$

## Dijkstra's Directly on Mesh Edges



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## 8x8 Cartesian Grid: Side Length 1



Shortest path along mesh is length $7 \sqrt{2}$

## Dijkstra's Directly on Mesh Edges

## 8x8 Cartesian Grid: Side Length 1



## Dijkstra's Directly on Mesh Edges

## 8x8 Cartesian Grid: Side Length 1



Shortest path along mesh is 14

## Dijkstra's Directly on Mesh Edges

Does refining the grid help?
$15 \times 15$ Cartesian Grid: Side Length 0.5


## Dijkstra's Directly on Mesh Edges

Does refining the grid help?
$15 \times 15$ Cartesian Grid: Side Length 0.5


Nope!

## Dijkstra's Directly on Mesh Edges

In general, mesh biases the solution!


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## Fast Marching

A modification of Dijkstra's algorithm to cut through triangles


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A modification of Dijkstra's algorithm to cut through triangles


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## Mesh Isomorphisms

An isomorphism preserves all pairwise geodesic distances


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## Mesh Isomorphisms

Contrast with Euclidean



